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The Liar Paradox: Considering Fuzzy Logic and Trivalent Truth Conditions

1 Introduction

The term *self-reference* is used to denote any set of circumstances in which someone or something refers to itself. In philosophy, self-reference is primarily studied in the context of language. A *self-referential sentence*, then, is a sentence which refers to itself, and perhaps the most notable instance of this is the *liar sentence*. Consider a sentence named (A), which says of itself (that is, says of (A)) that it is false:

(A) Sentence (A) is false.

(A) is self-referential because it refers to itself in the sentence. Additionally, it is *paradoxical* in the sense that it is self-contradictory. This self-contradictory nature of liar sentences is the primary motivator for what will follow in this paper.

In what follows I will propose two distinct solutions to the *liar paradox*, one that involves fuzzy logic and another that involves supplying a third truth value to bivalent logic. I will also present *revenge paradoxes* for these solutions. Revenge paradoxes are arguments that arise in response to some proposed solution to a paradox that prove that solution to be insufficient. I will show that both solutions suffer from revenge paradoxes and that thus neither solution is satisfactory when dealing with the liar paradox. First, I will explain why liar sentences invariably lead to a contradiction.

2 **Principle of Bivalence**

In "Function and Concept", Friedrich Ludwig Gottlob Frege offers his notion of a *function* (and, by extension, a *concept*) and demonstrates how he derives truth values of two types: true and false (Frege 137, 139). It is important to identify what suffices as truth values because we use them to assert propositions about given contexts or, more generally, the world. More specifically, truth values are important in the case of the liar paradox because we will only be able to discern the contradiction that emerges by evaluating the veracity of what is taken at face value from the liar sentence against the veracity of what is actually meant once one examines it more closely.

2.1 Functions

Frege defines a function as follows: "A function of *x* [is] taken to be a mathematical expression containing *x*, a formula containing the letter *x*. Thus...the expression $2x^3 + x$ would be a function of *x*, and $2[\times]2^3 + 2$ would be a function of 2." (Frege 131) He takes issue with the view that expressions such as "1 + 3" and "2 + 2" are equal but not the same because he believes that expressions, although able to vary in terms of form (that is, the way in which they are written), retain some intimate feature about them known as content. He called this content *Bedeutung*, and this roughly translates to "reference". Frege would state that, because both "1 + 3" and "2 + 2" are equal to four, both expressions retained the same *Bedeutung*. In the same vein, Frege supposes that expressions such as "2", "1 + 1", "3 – 1", and "6[/]3" all have the same *Bedeutung*. After all, there is no difference between the value, that being two, to which they all evaluate, but instead the only difference lies in their respective forms (Frege 132). Frege also shows how the *Bedeutung* of the word "function" is extended.

First, Frege asserts that "the field of mathematical operations", by which he means the ways in which we can manipulate numbers, has been extended (Frege 137). He writes, "Besides addition, multiplication, exponentiation, and their converses [subtraction, division, logarithms,

respectively], the various means of transition to the limit have been introduced" (Frege 137). Second, he asserts that the field of possible arguments and values for functions has been extended by the admission of complex numbers (Frege 137). With respect to both these directions, Frege began adding to the signs (+), (-), etc., so that he could better construct functional expressions, and he also added signs such as (=), (>), (<), etc., so that he could evaluate functions where *x* takes the place of an argument that is compared to some alternative expression (for example, the function $x^2 = 1$) (Frege 137). All this leads us to Frege's notion of a concept.

2.2 Concepts

Frege defined a concept as follows: "a concept is a function whose value is always a truth-value." (Frege 139) But where does he recover what he considers "truth values" in the first place? Frege has us consider the function $x^2 = 1$, where *x* takes the place of some numeric argument, and evaluates it for different arguments. He writes:

Now if we replace *x* successfully by -1, 0, 1, 2, we get:

$$(-1)^2 = 1,$$

 $0^2 = 1,$
 $1^2 = 1,$
 $2^2 = 1.$

Of these equations the first and third are true, the others false. (Frege 137)

From this Frege shows that the value of a function can be a truth value and, if so, that such truth values come in two varieties: true and false. *True functions* are functions whose resultant values are consistent with what is asserted by their argument values, whereas *false functions* are functions whose resultant values fail to be consistent in this manner. Thus, equations such as " 2^2

= 4", for example, are true because two raised to the second power does, in fact, equate to four, whereas equations such as " $2^2 = 1$ " are false because two raised to the second power does not equate to one. And as I stated before, Frege extended his terminology to more than just equations, and the aforementioned truth values also appear applicable in these cases. He writes:

Accordingly,

$$2^{2} = 4^{2}, 2 > 1^{2}, 2^{4} = 4^{2},$$

all stand for the same thing...the True, so that in

$$(2^2 = 4) = (2 > 1)$$

we have a correct equation. (Frege 137)

One might object by suggesting that the expressions " $2^2 = 4$ " and "2 > 1" cannot be equated because both express different claims. However, based on what we know about truth values, we can determine that " $(2^2 = 4) = (2 > 1)$ " may be reduced to "True = True", and in this way they express the very same claim and thus have the same *Bedeutung*.

3 Liar Paradox

Recall sentence (A) from Section 1:

(A) Sentence (A) is false.

The *principle of bivalence* asserts that every sentence expressing a proposition has exactly one truth value, either true or false, and I have shown in Section 2 how Frege arrived at this notion. Though, the difficulty in applying this principle with regard to (A) arises from the fact that a contradiction emerges under both the supposition that (A) is true and the supposition that (A) is false.

Suppose that (A) is true. If (A) is true, then what (A) asserts must be true, and (A) asserts that (A) is false. If (A) is determined to be false, then it certainly cannot be true. So by supposing

that (A) is true, we arrive at the conclusion that (A) is false, a contradiction. The converse can also be established.

Suppose that (A) is false. If (A) is false, then what (A) asserts must be false, and (A) asserts that (A) is false. In a case such as this, where the assertion that (A) is false is, in fact, false, then what (A) asserts of itself must actually be true. So by supposing that (A) is false, we arrive at the conclusion that (A) is true, again a contradiction.

Clearly, when we suppose that (A) is either true or false, we find that it ends up being both true and false. This is a paradox because propositions may be either true or false, but not both, under our current understanding. So how can we resolve this?

4 Fuzzy Logic

Fuzzy logic is a form of many-valued logic in which the truth value of propositions may be defined by any real number between 0 and 1 (inclusive). It is employed to handle cases of partial truth, where the truth value may range between being entirely true and entirely false. For example, consider the temperature of tap water. In Boolean logic, where the truth values of propositions may only be defined by the integer values 0 (representing false) or 1 (representing true), we would represent the temperature of tap water as either hot or cold. With fuzzy logic, we could instead represent the temperature by means of a gradient from hot to cold. That is, we could account for arbitrary ranges of the temperature that fall between being hot and cold with terms such as "slightly hot", "lukewarm", "slightly cold", etc., instead of being limited only to the prior two terms. As can be inferred, fuzzy logic is meant to be used to model logical reasoning in situations where the veracity of propositions is vague or imprecise, such as when the constraints of a particular proposition suggest that there are more outcomes to consider than just that of being entirely true or entirely false. Fuzzy logic is then a part of a family of many-valued

logics where truth values are interpreted as degrees of truth instead of in accordance with the principle of bivalence.

4.1 **Definitions**

Fuzzy logic works with membership values in a way that mimics Boolean logic, and Lotfi Zadeh's 1965 proposal of fuzzy set theory suggests that we should be able to determine the *complement* of a fuzzy set by relating it to Boolean logic (Zadeh 340). Under Boolean logic, the complement can be expressed as follows: NOT(x). We will need to concern ourselves with a fuzzy replacement for this.

The most important thing to realize about fuzzy logic is that it is a superset of conventional Boolean logic. That is, if we set the fuzzy logic values at their extremes of 1 (entirely true) and 0 (entirely false), then we will find that the standard logical operations from Boolean logic hold. Below I have constructed the standard truth table for the NOT operator (see table 1).

Table 1

Truth Table for NOT Operator Under Boolean Logic

Α	not A
0	1
1	0

Considering that, under fuzzy logic, the truth of any proposition is a matter of degree, we will need to reframe this truth table such that it accounts for the range of possible degrees of truth that extend from 0 to 1. We can do this by replacing the operation NOT *A* with the operation 1 - A. We will find that the previous truth table goes unaffected by this replacement (see table 2).

Table 2

T 41 T	11 0	NIOT	\mathbf{O}	TT 1	г	т .
I mith I	able for		Operator	Under	FIIZZV	LOGIC
II addi I		1101	operator	Onder	Iuzzy	LUGIU

Α	1 – A
0	1
1	0

Thus, 1 - A suffices as the complement for A under fuzzy logic.

4.2 Initial Solution

To successfully evaluate the liar sentence under fuzzy logic, we will need to interpret it under a more formal construction than how I have presented it thus far. Let x denote the degree of truth of sentence (A). If we want to determine x, then we must construct (A) in such a way that the choice of its possible truth values are restricted. I have shown that, so far, the only possible truth values that a sentence may evaluate to are true and false, so these will comprise the scope of truth values to which (A) will adhere.

As shown in Section 2, Frege's conception of truth has it that sentences are referring terms. That is, they are names that name what is true and what is false. So one might suggest, if we had a Fregean version of fuzzy logic, that we might say that sentences are names of degrees of truth between 0 and 1. For example, it might make sense to say something like "(1 + 3 = 4) = 1" because "1 + 3 = 4" is just a name for the degree of truth 1, and so is "1". Recall sentence (A) from Section 1:

(A) Sentence (A) is false.

On the view that sentences are names of degrees of truth, it makes sense to say, for example, "Sentence (A) is false = 0.5". After all, both sides of the equation contain a name of a degree of truth. However, it still will not make sense to say "(A) = 0.5". Under Frege's conception of truth, the left-hand side of the equation is a name of a *sentence*, not the name of a *degree of truth*. In other words, it is a name of a *name of a degree of truth*. If we want to use names of sentences to say something about degrees of truth, then we will need some device that allows us to take the names of sentences as arguments and map them onto degrees of truth.

Suppose that we have a degree function deg(x) that takes names of arguments and maps them onto degrees of truth. Because we want to evaluate (A) using this function, we end up with the degree function deg("Sentence (A) is false.") or, more simply, deg("False(A)"). Under Boolean logic, supposing that some proposition is false is equivalent to invoking the NOT operator onto it. Thus, we end up with the following equation:

$$deg(A) = NOT(deg(A))$$

We can generalize the NOT operator to its equivalent under fuzzy logic:

$$\deg(\mathbf{A}) = 1 - \deg(\mathbf{A})$$

We can rearrange this equation so that both deg(A)s are on the same side of the equation:

$$2(\deg(A)) = 1$$

We can divide both sides of the equation by 2 to solve for deg(A):

$$deg(A) = 0.5$$

So there appears to be a non-paradoxical solution to the liar paradox under fuzzy logic. Instead of treating the liar sentence as both entirely true and entirely false, we can treat it as precisely half-true and half-false.

4.3 Background to Revenge Paradox

In "On Sinn and Bedeutung", Frege argues that every word, every expression, and every sentence has two semantic values: *sense* and *reference* (Frege 155). He asserts that both sense and reference are compositional, meaning that the sense of a compound structure (in this case, a sentence) is a function of the senses of its component parts, and likewise for reference. Frege

writes, "A proper name (word, sign, combination of signs, expression) *expresses* its sense, *stands for* [*bedeutet*] or *designates* [*bezeichnet*] its *Bedeutung*. By employing a sign we express its sense and designate its *Bedeutung*." (Frege 156) To put it more clearly, the senses which sentences express are propositions, whereas their referents are the truth values (which, in turn, constitute the *Bedeutung*). And again, Frege asserts that the truth values may only be either true or false (Frege 157).

Because the reference of a sentence is a function of the reference of the parts of the sentence, we should be able to extract subsentential (that is, less than a sentence) terms out of the sentence and substitute them with terms which have the same reference without causing any harm to (that is, changing) the truth value of the sentence (Frege 166). Consider the following sentence:

The Morning Star is a body illuminated by the Sun. (Frege 156) This sentence is true because "the Morning Star" is another name for the planet Venus, and Venus is, in fact, illuminated by the Sun. So since this sentence is true, its referent, according to Frege, is the True. Now, the term "the Morning Star" happens to refer to the same thing as another term, "the Evening Star". And again, because the reference of a sentence is a function of the reference of its parts, we should be able to swap out terms for coreferential ones without changing the truth value of the sentence. Thus, the following sentence should also be true, since the referent has not been changed:

The Evening Star is a body illuminated by the Sun. (Frege 156) And this sentence is also true. Now, Frege is sympathetic to the idea that someone might determine that one of the two sentences is true and that the other is false (Frege 156). Such an outcome would make sense if the person in question did not know that both the Morning Star

and the Evening Star referred to the same thing. However, this misconception would have no bearing on the *Bedeutung* of the sentence, since again the referent was not changed. Thus, the thought corresponding to each sentence cannot be its reference, but rather it must be considered what Frege called its sense (Frege 156). In "Letter to Husserl, 24.5.1891", Frege writes, "Judgment in the narrower sense could be characterized as a transition from a thought to a truth-value." (Frege 150) And he believes that such a schema could look like this:

concept word

 \downarrow

sense of the concept word

(sense)

 \downarrow

object falling under the concept (Frege 150)

Frege would state that this schema similarly holds true for proper names and sentences, with the only difference between the three being their scopes: concept words usually can refer to more than one thing, proper names refer to just one thing, and sentences refer to just the sentence under consideration (Frege 150). Considering all this, is there an issue which emerges from the move that I made earlier (hereinafter referred to as the "fuzzy liar" and abbreviated as "FL")?

4.4 **Revenge Paradox**

As we all must recognize by now, fuzzy logic allows us to deal with many more truth values than Frege had ever anticipated existing. However, no issue should arise when we combine this framework with Frege's principle that substituting terms for coreferential ones should never change the truth value of a sentence. Let us, then, consider two cases under the supposition that the fuzzy liar is true to some degree that is less than 1: one where the truth value of the fuzzy liar is (1) some degree that is less than 1 and another where the truth value is (2) equal to 1. It will be clear why the two cases result in very different problems.

In the first case, the fuzzy liar is true to some degree that is less than 1. More formally, we can express this as follows:

(FL) deg(FL) < 1.

Suppose that the degree of truth of (FL) is equal to 0.5. Under this supposition, "deg(FL)" and "0.5" are coreferential terms. If Frege's principle of substitution holds true, then we should be able to substitute terms for coreferential ones in (FL) without changing the truth value of (FL). That is, the following sentence should retain the same truth value as (FL):

(FL*) 0.5 < 1.

This sentence states that 0.5 is less than 1. It should be obvious that this is entirely true, and, under fuzzy logic, the truth value of something that is entirely true is 1. However, since we originally supposed that the degree of truth of (FL) was equal to 0.5, then this would suggest that this very same assertion is instead only half-true. Thus, a contradiction emerges where we are somehow equating 0.5 with 1. To put it more clearly, notice that, based on our supposition, all of the following will have to be coreferential:

"deg(FL)", "0.5", "deg(FL*)", "1"

So by assigning the degree of truth 0.5 to (FL), we commit to the mathematical absurdity that 0.5 is equal to 1. And a similar argument will hold for any degree of truth of (FL) that is less than 1. For example, if we assign the degree of truth 0 to (FL), then we will arrive at the conclusion that 0 = 1, and this is obviously false. And the same would hold true for the degree of truth 0.99 if it were assigned to (FL): we will arrive at the conclusion that 0.99 = 1, and again this is obviously false.

In the second case, the fuzzy liar is, again, true to some degree that is less than 1:

(FL) deg(FL) < 1.

However, suppose that the degree of truth of (FL) is equal to 1 instead. Under this supposition, "deg(FL)" and "1" are coreferential terms. Again, if Frege's principle of substitution holds true, then we should be able to substitute terms for coreferential ones in (FL) without changing the truth value of (FL). That is, the following sentence should retain the same truth value as (FL):

Under fuzzy logic, something that has a truth value of 1 is entirely true. However, this sentence states that 1 is less than itself, and this cannot be true under any set of circumstances. In this case, then, this sentence must be entirely false. Thus, the degree of truth of (FL*) ought to be 0; not 1. So by assigning the degree of truth 1 to (FL), we are substituting a term for an allegedly coreferential one and thus altering the truth value of the sentence (unless 0 = 1 or 1 < 1, both of which are false). It seems as though we are left wondering if there is any way that we can frame this issue such that it does not lead to mathematical absurdity.

4.5 Assessment

Under the original liar paradox, fuzzy logic allows us to assert that the truth value of the liar sentence is 0.5. This would be an acceptable solution, since it resolves the contradiction of the liar sentence being both entirely true and entirely false by instead showing that it is half-true and half-false, had a revenge paradox not emerged from it. The revenge paradox I presented permits us to assert logically impossible conclusions, such as recognizably different values somehow equating with each other or the exact same value somehow equating with a value that is less than itself, so it does not seem as though we can actually establish a valid solution to the liar paradox under this approach.

5 Trivalent Truth Conditions

Another way that we might attempt to resolve the liar paradox involves introducing a third truth value, one which accounts for sentences that might best be described as *indeterminate*, a state of truth in which the veracity of some proposition cannot be precisely determined. Of course, justification is necessary for something like this.

5.1 Background

In the chapter "Three-valued logic: beginnings" of *Vagueness*, Timothy Williamson provides a reasonable motive for a system of trivalent logic:

[C.S. Peirce's] idea was that two-valued logic, although not wholly incorrect, is valid only within a limited domain, [trivalent logic] being needed for full generality...[trivalent logic] is that logic which, though not rejecting entirely the Principle of Excluded Middle, nevertheless recognizes that every proposition, S is P, is either true, or false, or else S has a lower mode of being such that it can neither be determinately P, nor determinately not P, but is at the limit between P and P'. (Williamson 102)

One might suppose that Peirce was concerned with vagueness because, as we already know by now, bivalent logic just does not provide any answers with regard to cases of partial truth. However, a different, more pragmatic justification for trivalent logic was later offered by Jan Łukasiewicz, who was concerned with free will (Williamson 102). Williamson writes, "He thought that fatalism can be avoided only if some statements about the future, such as 'There will be a sea-fight tomorrow', are not yet true or false." (Williamson 102) This rationale should be intuitively pleasing to us. After all, how would we possibly be able to determine the truth value of something that has not yet occurred?

Now, the *principle of excluded middle* states that either a proposition or its negation is true. But if propositions are permitted to exist in the aforementioned "lower mode" where they may neither be considered determinately true nor determinately false, then one might question how this system of trivalent logic does not entirely reject the principle in question.

If we take into consideration the idea that something that is not yet true or false must mean that it is not yet determined, then it seems reasonable to withhold probability judgments across any of the possible outcomes of a proposition when we are unable to determine its truth value. Accordingly, below I have constructed some truth tables for negation, conjunction, disjunction, the material conditional, and the material biconditional under trivalent logic, where "T" stands for true, "F" stands for false, and "#" stands for indeterminate (see tables 3, 4, 5, 6, and 7).

Table 3

Truth Table for Negation Under Trivalent Logic

Р	¬₽
Т	F
#	#
F	Т

When P is assigned the truth value #, the truth value of \neg P is also #. In reference to the sea-fight example, if we cannot determine that there will be a sea-fight tomorrow, then we certainly cannot determine that there will *not* be a sea-fight tomorrow.

Table 4

Truth Table for Conjunction Under Trivalent Logic

Р	Q	$\mathbf{P} \wedge \mathbf{Q}$
Т	Т	Т
Т	#	#
Т	F	F
#	Т	#
#	#	#
#	F	F
F	Т	F
F	#	F
F	F	F

When P is assigned the truth value T and Q is assigned the truth value #, the truth value of P \land Q is #; for a conjunction to be true, both propositions have to be true. When both P and Q are assigned the truth value #, the truth value of P \land Q is #; both propositions are indeterminate, so that is the only conclusion that we can derive. When P is assigned the truth value F and Q is assigned the truth value #, the truth value of P \land Q is F because a conjunction is false if at least one of its conjuncts is false.

Table 5

Р	Q	$\mathbf{P} \lor \mathbf{Q}$
Т	Т	Т
Т	#	Т
Т	F	Т
#	Т	Т

Truth Table for Disjunction Under Trivalent Logic

#	#	#
#	F	#
F	Т	Т
F	#	#
F	F	F

When P is assigned the truth value T and Q is assigned the truth value #, the truth value of P \lor Q is T because the truth of one of the disjuncts is sufficient to guarantee the truth of the entire disjunction. When both P and Q are assigned the truth value #, the truth value of P \lor Q is #; both propositions are indeterminate, so that is the only conclusion that we can derive. When P is assigned the truth value F and Q is assigned the truth value #, the truth value of P \lor Q is # because, again, if one of the disjuncts is true, then the entire disjunction is true. We would just not yet be in a position to know if the indeterminate disjunct were true.

Table 6

Р	Q	$P \rightarrow Q$
Т	Т	Т
Т	#	#
Т	F	F
#	Т	Т
#	#	#
#	F	#
F	Т	Т
F	#	Т

Truth Table for Material Conditional Under Trivalent Logic

F	F	Т

When P is assigned the truth value T and Q is assigned the truth value #, the truth value of $P \rightarrow Q$ is # because when the first proposition is true, the truth value of the material conditional is the same as the truth value of the second proposition. When P is assigned the truth value # and Q is assigned the truth value T, the truth value of $P \rightarrow Q$ is T because the truth value for every material conditional where the second proposition is true is T. When both P and Q are assigned the truth value #, the truth value of $P \rightarrow Q$ is #; both propositions are indeterminate, so that is the only conclusion that we can derive. When P is assigned the truth value # and Q is assigned the truth value of $P \rightarrow Q$ is # because the truth value # and Q is assigned the truth value F, the truth value of $P \rightarrow Q$ is # because the truth value of T or F. When P is assigned the truth value F and Q is assigned the truth value #, the truth value F and Q is assigned the truth value #, the truth value F and Q is assigned the truth value #, the truth value F and Q is assigned the truth value # and Q is T because the truth value of $P \rightarrow Q$ is T because the truth value of $P \rightarrow Q$ is T because the truth value F and Q is assigned the truth value #, the truth value of $P \rightarrow Q$ is T because the truth value of $P \rightarrow Q$ is T because the truth value of $P \rightarrow Q$ is T because the truth value of $P \rightarrow Q$ is T because the truth value of $P \rightarrow Q$ is T because the truth value of $P \rightarrow Q$ is T because the truth value of $P \rightarrow Q$ is T because the truth value of $P \rightarrow Q$ is T because the truth value of $P \rightarrow Q$ will be T if Q has either the truth value T or F.

Table 7

Truth Table for Material Biconditional Under Trivalent Logic

Р	Q	$\mathbf{P} \leftrightarrow \mathbf{Q}$
Т	Т	Т
Т	#	#
Т	F	F
#	Т	#
#	#	Т
#	F	#
F	Т	F
F	#	#

F	F	Т

Because we have the truth tables of both the conjunction and the material conditional under trivalent logic, we can derive the truth table of the material biconditional. Here, propositions have the truth value T when they share the same truth value and the truth value F when the truth values are different. Note that when only one of the propositions has the truth value # the truth value of $P \leftrightarrow Q$ is also #. This is because that indeterminate proposition *could* end up sharing the same truth value as the determinately true/false proposition, but we would just not yet be at a point where we could determine that.

So it seems as though a system of trivalent logic does not entirely reject the principle of excluded middle because a truth value with any determinate meaning to it (so not denoted by "#") is only assigned to a proposition once we know that the truth value is either determinately true or determinately false.

5.2 **Definitions**

In "Making Sense of (in)Determinate Truth: The Semantics of Free Variables", John Cantwell provides definitions which allow us to understand indeterminateness more formally:

(1) $A_1, ..., A_n$ semantically entail *B* iff *B* is true relative to every assignment *g* where all the A_i are true.

(2) In a context where all has been assumed is $A_1, ..., A_n$, an assignment g is *admissible* if and only if each A_i is true relative to g.

(3) A sentence *A* is *determinately true (false)* in a context iff *A* is true (false) in every admissible assignment in that context.

(4) A sentence *A* has *indeterminate truth value* in a context iff *A* is neither determinately true nor determinately false in that context. (Cantwell 2723–2724)

We can establish two conclusions from these definitions. First, while assignment succeeds when some general solution includes *at least one* possible solution to some particular function, entailment only succeeds when some general solution includes *all* possible solutions to it. Second, because some sentence *x* is determinately true or determinately false only if it accounts for *all* possible solutions given a particular context, the sentence is indeterminate if it fails to account for *at least one* possible solution.

5.3 Initial Solution

Recall sentence (A) from Section 1:

(A) Sentence (A) is false.

To show that this sentence would be more appropriately described as indeterminate, I will need to demonstrate that the truth value of this proposition never ends up being only true or only false.

For "Sentence (A) is false" to be determinately true or determinately false, then it must be either true or false in every admissible assignment in that context, that is to say it must be only true or only false in all possible cases. We already know from Section 3 that if we suppose that (A) is true, then (A) will end up being false, and that if we suppose that (A) is false, then (A) will end up being false, and that if we suppose that (A) is false, then (A) will end up being true. The fact that the sentence always ends up being true or false does *not* make it both determinately true and determinately false. Again, for something to be determinately true or determinately false, only *one* of the truth values may hold for all cases. "Sentence (A) is false" is not true or false in all possible cases, so it is not determinately true or determinately false. And if "Sentence (A) is false" is not determinately true or determinately false, then it must retain an indeterminate truth value.

5.4 Revenge Paradox

Consider a variation of sentence (A) where the sentence is *not true* instead of false:

(A*) Sentence (A*) is not true.

One might initially wonder what exactly differentiates sentence (A) from sentence (A*). Put simply, sentence (A*)'s scope of truth values is more broad than sentence (A)'s because it not only captures that the sentence under investigation is false, but that the sentence is *either* false or indeterminate.

Suppose that (A^*) is true. If (A^*) is true, then what (A^*) asserts must be true, and (A^*) asserts that (A^*) is not true, meaning that (A^*) is either false or indeterminate. If (A^*) is determined to be either false or indeterminate, then it certainly cannot be true. So by supposing that (A^*) is true, we arrive at the conclusion that (A^*) is either false or indeterminate, a contradiction.

Now, suppose that (A^*) is false. If (A^*) is false, then what (A^*) asserts must be false, and (A^*) asserts that (A^*) is not true, again meaning that (A^*) is either false or indeterminate. In a case such as this, where the assertion that (A^*) is not true is, in fact, not true, then what (A^*) asserts of itself must actually be true. So by supposing that (A^*) is false, we arrive at the conclusion that (A^*) is true, again a contradiction. And this line of reasoning holds true under the supposition that sentence (A^*) is indeterminate as well.

Suppose that (A^*) is indeterminate. If (A^*) is indeterminate, then what (A^*) asserts must be indeterminate, and (A^*) asserts that (A^*) is not true, again meaning that (A^*) is either false or indeterminate. In a case such as this, where the assertion that (A^*) is not true is, in fact, not true, then what (A^*) asserts of itself must actually be true.

It seems as though we are back in a position where all our suppositions lead to a contradiction. This suggests that introducing further truth values is unlikely to aid us in resolving the liar paradox.

5.5 Assessment

Under the original liar paradox, trivalent logic allows us to assert that the truth value of the liar sentence is indeterminate. If this holds true, then the liar sentence cannot be said to be either determinately true or determinately false. So instead of us having to deal with the contradiction that arises from the liar sentence being both entirely true and entirely false, we could withhold judgment and not have to concern ourselves with the contradiction at all. However, the revenge paradox I presented has it that the liar sentence is more broad to the point such that its being indeterminate even leads to a contradiction, so, like with fuzzy logic, it does not seem as though we can actually establish a valid solution to the liar paradox under this approach either.

6 Conclusion

In this paper I have shown that the liar paradox seemingly prevails under all three systems of bivalent logic, fuzzy logic, and trivalent logic. It is evident throughout the paper that bivalent logic simply fails to resolve the contradiction entailed by the liar paradox, whereas both fuzzy logic and trivalent logic succumb to revenge paradoxes which ultimately invalidate them as proper solutions. That is, none of the approaches seem to sufficiently resolve the liar paradox.

However, if I had to decide which approach seemed the *most* satisfactory considering this overall failure, then I would say that fuzzy logic does. Under trivalent logic, we go through so much to introduce an entirely new truth value just to end up in a similar situation to that of where we were under bivalent logic, with the only difference being that our issue now extends an additional truth value. With the initial solution involving fuzzy logic, we at least end up deriving

a truth value for the original liar paradox which makes intuitive sense, but still we ultimately end up with mathematical absurdity. Perhaps there is a way to frame the issue that the revenge paradox sheds light on that I am simply not aware of at this point in time.

The lesson to be found here does not seem to be that liar sentences are neither true nor false, but instead that, in order to reflect as many of our intuitions as possible, paradoxical sentences must somehow be permitted without our theory of truth failing. And while it seems as though discarding bivalence provides us with one way of doing this, it is still unclear what, exactly, is the right way to approach liar sentences.

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